

# Solvers Principles and Architecture (SPA)

## *General Introduction*

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Before understanding **Solvers**

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We need to talk about **Problems**

# What makes a problem important ?

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## Travelling Salesman Problem

Given a list of cities and the distances between each pair of cities, what is the shortest possible route that visits each city exactly once and returns to the origin city?

# What makes a problem important ?

Reduction of other problems.

$$\begin{pmatrix} \textit{Problem1} \\ \vdots \\ \textit{ProblemN} \end{pmatrix} \rightsquigarrow \textit{ProblemA}$$

- The transformation (reduction) may be non-trivial to find
- Needs to be “**simpler**” than solving the new problem



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Reduction example

## Satisfiability

Given a set  $V$  of Boolean variables and a collection  $C$  of clauses over  $V$ , is there a satisfying truth assignment for  $C$ ?

## Quadratic Diophantine Equations

Given positive integers  $a$ ,  $b$ , and  $c$ , are there positive integers  $x$  and  $y$  such that  $ax^2 + by^2 = c$ ? (Transformation from 3SAT [Manders and Adleman 1978].)

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Poincaré said so!

- Its **long resistance** (beyond the current state-of-the art methods)
- Requires **new insights** (connections, perspectives) to get solved
- Example: Hilbert's famous list of problems (1900)
- Example: Millennium Prize Problems

## Riemann Hypothesis

All the non-trivial zeros of the Riemann zeta function have their real part equal to  $\frac{1}{2}$ .

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- 1 What Makes a Problem Important?
- 2 Important Problems**
- 3 Solving in Mathematics
- 4 Solving Computer Science
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## Satisfiability (DPLL algorithm)

Is there a Boolean assignment that satisfies

$$(v_1 \vee \bar{v}_2) \wedge (\bar{v}_1 \vee v_2)$$

## Quantifier Elimination (Cylindrical Algebraic Decomposition)

Is the following sentence true over the reals

$$\forall a, b. \exists x. \quad x^2 + ax + b = 0$$

# Important Problems You Must Be Aware Of

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## Convex Optimization (SemiDefinite Programming)

$$\begin{aligned} \text{Min/Max} \quad & C \bullet X \\ \text{Subject to} \quad & A_i \bullet X = b_i, \quad i = 1, \dots, m \\ & X \succeq 0 \end{aligned}$$

## Differential Equations (Numerical Algorithms)

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**Object of study** in mathematics is the **set of solutions** of equations

$$f(x) = 0$$

## Nature and operators in $f$

- Linear (vector of)
- Polynomial (vector of)
- With special operators: derivations
- ...

## Solution Space

- Finite fields ( $\mathbb{Z}/p\mathbb{Z}$ )
- Reals
- Differential functions
- Probability densities
- ...



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- Exact Solving (exact general form of the solution)
- Properties of the set of solutions (finiteness, bounded, structure, symmetries etc.)
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Purpose: **Classification**

- 1 What Makes a Problem Important?
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## Definition

- Description of the parameters
- Statement of what an answer, or solution, is required to satisfy

## Example: Traveling Salesman Problem

The problem consists of a finite set of locations/cities  $C = \{c_1, \dots, c_m\}$  and for each pair of cities  $c_i, c_j$  in  $C$ , the distance  $d(c_i, c_j)$  between them. A solution is an ordering  $c_{\pi(1)}, c_{\pi(2)}, \dots, c_{\pi(m)}$  such that minimizes

$$\left( \sum_{i=1}^{m-1} d(c_{\pi(i)}, c_{\pi(i+1)}) \right) + d(c_{\pi(m)}, c_{\pi(1)})$$

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**Step-by-step** procedure to solve **any** instance of a given problem.

## Efficiency

- Time complexity (is not the only important parameter)
- How does the **time** needed to solve the problem evolves when the **input length** increases?
- Number of symbols in the description of the instance with respect to the encoding scheme for the problem.

## Example of input length

- Alphabet  $\{c, [, ], /, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$
- "c[1]c[2]c[3]c[4]//10/5/9//6/9//3" (32 symbols)

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## Polynomial Time Complexity

Time complexity is  $O(p(n))$  for some polynomial  $p$  with input length  $n$ .

## Exponential Time Complexity

Time complexity cannot be bounded by a polynomial time complexity (including  $n^{\log n}$ ).

	$n = 10$	$n = 30$	$n = 60$
$n^3$	0.001 s	0.027 s	0.216 s
$3^n$	0.059 s	6.5 years	$1.3 \times 10^{13}$ centuries

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- Complexity (worst case with respect to the input length)
- Target special cases
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- Dive into three different problems: **SMT**, **CAD**, **SDP**
- Study their related solvers: **Algorithms** and **Data Structure**
- Learn how to **Deconstruct** a problem



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