

TD2

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Natural transformations

Exercise 0 Let $\mathbf{2}$ be the discrete category with two objects. Prove that the functor category $[\mathbf{2}, \mathcal{B}]$ is isomorphic to the product category $\mathcal{B} \times \mathcal{B}$. (The notation \mathcal{B}^2 for $[\mathbf{2}, \mathcal{B}]$ is now justified.)

Exercise 1 Let \mathcal{A} and \mathcal{B} be categories. Prove that $[\mathcal{A}, \mathcal{B}]^{\text{op}} \cong [\mathcal{A}^{\text{op}}, \mathcal{B}^{\text{op}}]$.

Exercise 2 Let $\mathcal{A} \begin{array}{c} \xrightarrow{F} \\ \Downarrow \alpha \\ \xrightarrow{G} \end{array} \mathcal{B}$ be a natural transformation. Prove that α is a natural isomorphism if and only if its components α_A are isomorphisms (in \mathcal{B}) for each $A \in \mathcal{A}$. (This can be taken as an equivalent definition to natural isomorphism.)

Exercise 3 A permutation of a set X is a bijection $X \rightarrow X$. Let $\mathbf{Sym}(X)$ denote the set of permutations of X . A total order on a set X is an order \leq such that for all $x, y \in X$, either $x \leq y$ or $y \leq x$. So a total order amounts to a way to placing the elements of X in a sequence. Let $\mathbf{Ord}(X)$ denote the set of total orders of X . Let \mathcal{A} denote the category of finite sets and bijections.

- (a) There is a canonical way to regard both \mathbf{Sym} and \mathbf{Ord} as functors from \mathcal{A} to the category of sets \mathbf{Set} . Make it explicit.
- (b) Show that there is no natural transformation $\mathbf{Sym} \rightarrow \mathbf{Ord}$.
- (c) For an n -element set X , how many elements do the sets $\mathbf{Sym}(X)$ and $\mathbf{Ord}(X)$ have?

Conclude that $\mathbf{Sym}(X) \cong \mathbf{Ord}(X)$ but not naturally for all $X \in \mathcal{A}$. (This is an example where maps in \mathcal{A} are isomorphic in the standard sense with no natural way to match them up.)

Exercise 4 A functor $F : \mathcal{A} \rightarrow \mathcal{B}$ is *essentially surjective on objects* if for all $B \in \mathcal{B}$, there exists $A \in \mathcal{A}$ such that $F(A) \cong B$. Prove that a functor is an equivalence if and only if it is faithful, full, and essentially surjective on objects.

Exercise 5 (bonus) Show that equivalence of categories is an equivalence relation.